Java Atomic Classes & Operations: Introduction



Douglas C. Schmidt

<u>d.schmidt@vanderbilt.edu</u>

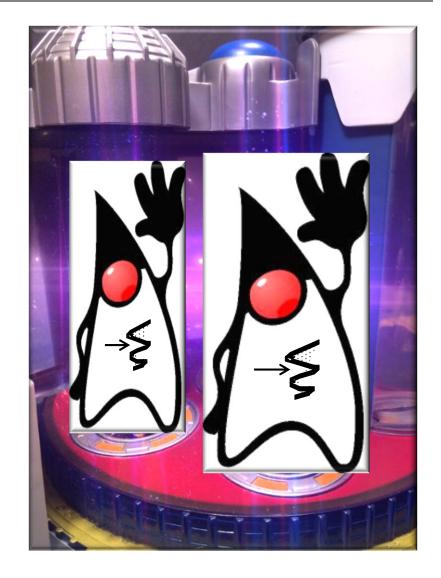
www.dre.vanderbilt.edu/~schmidt

Institute for Software Integrated Systems Vanderbilt University Nashville, Tennessee, USA



Learning Objectives in this Part of the Lesson

 Understand how Java atomic classes & operations provide concurrent programs with lock-free, thread-safe mechanisms to read from & write to single variables



Learning Objectives in this Part of the Lesson

- Understand how Java atomic classes & operations provide concurrent programs with lock-free, thread-safe mechanisms to read from & write to single variables
- Note a human known use of atomic operations



 The java.util.concurrent.atomic package several types of atomic actions on objects

Package java.util.concurrent.atomic

A small toolkit of classes that support lock-free thread-safe programming on single variables.

See: Description

Class Summary	
Class	Description
AtomicBoolean	A boolean value that may be updated atomically.
AtomicInteger	An int value that may be updated atomically.
AtomicIntegerArray	An int array in which elements may be updated atomically.
AtomicIntegerFieldUpdater <t></t>	A reflection-based utility that enables atomic updates to designated volatile int fields of designated classes.
AtomicLong	A long value that may be updated atomically.
AtomicLongArray	A long array in which elements may be updated atomically.
AtomicLongFieldUpdater <t></t>	A reflection-based utility that enables atomic updates to designated volatile long fields of

See docs.oracle.com/javase/8/docs/api/java/util/concurrent/atomic/package-summary.html

- The java.util.concurrent.atomic package several types of atomic actions on objects
 - Atomic variables
 - Provide lock-free & thread-safe operations on single variables



See docs.oracle.com/javase/tutorial/essential/concurrency/atomicvars.html

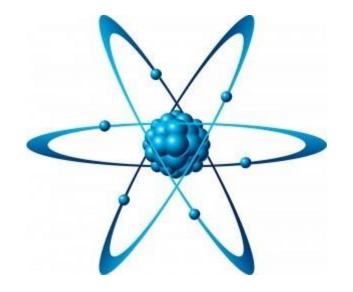
- The java.util.concurrent.atomic package several types of atomic actions on objects
 - Atomic variables
 - Provide lock-free & thread-safe operations on single variables
 - e.g., AtomicLong supports atomic "compare-and-swap" operations

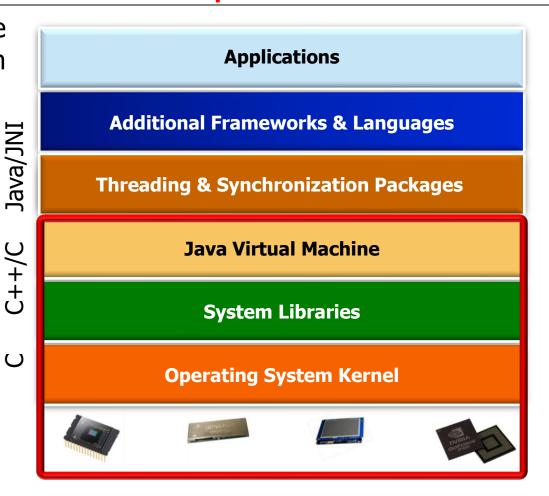
```
<<.lava Class>>
               AtomicLong
AtomicLong(long)
AtomicLong()
fget():long
set(long):void
lazySet(long):void
getAndSet(long):long
compareAndSet(long,long):boolean
weakCompareAndSet(long,long):boolean
getAndIncrement():long
getAndDecrement():long
getAndAdd(long):long
fincrementAndGet():long
decrementAndGet():long
addAndGet(long):long
getAndUpdate(LongUnaryOperator):long
updateAndGet(LongUnaryOperator):long
getAndAccumulate(long,LongBinaryOperator):long
accumulateAndGet(long,LongBinaryOperator):long
toString()
intValue():int
longValue():long
floatValue():float
doubleValue():double
```

- The java.util.concurrent.atomic package several types of atomic actions on objects
 - Atomic variables
 - LongAdder
 - Allows multiple threads to update a common sum efficiently under high contention

- add(long):void
- increment():void
- decrement():void
- sum():long
- reset():void
- sumThenReset():long
- toString()
- longValue():long
- intValue():int
- floatValue():float
- doubleValue():double

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void lock(int *mutex) {

continue;

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int compareAndSwap(int *loc,
                                             int expected,
                                             int updated) {
                           START ATOMIC();
                           int oldValue = *loc;
                           if (oldValue == expected)
                              *loc = updated;
                           END ATOMIC();
                           return oldValue;
while (compareAndSwap(mutex, 0, 1) == 1)
```

The lock() method uses compareAndSwap() to implement mutual exclusion (mutex) via a "spin-lock"

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compareAndSwap() must be called only once per lock attempt

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```

If compareAndSwap() returns 1 that means the mutex is "acquired" so the loop keeps spinning

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continue;

START ATOMIC();

*mutex = 0;

END ATOMIC();

```
int compareAndSwap(int *loc,
                                                 int expected,
                                                 int updated) {
                               START ATOMIC();
                               int oldValue = *loc;
                               if (oldValue == expected)
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                               END ATOMIC();
                               return oldValue:
void lock(int *mutex) {
  while (compareAndSwap(mutex, 0, 1) == 1)
void unlock(int *mutex) {
                                 The unlock() method atomically
                                   resets the mutex value to 0
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```
void unlock(int *mutex) {
  START ATOMIC();
  *mutex = 0;
  END ATOMIC();
```

continue;

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                             END ATOMIC();
                             return oldValue:
while (compareAndSwap(mutex, 0, 1) == 1)
                               The unlock() method atomically
                                 resets the mutex value to 0
```

```
void unlock(int *mutex) {
  START ATOMIC();
  *mutex = 0;
  END ATOMIC();
```

- Atomic operations can be implemented other ways
 - e.g., "test-and-set"

```
int testAndSet(int *loc) {
   int oldValue;
   START_ATOMIC();
   oldValue = *loc;
   *loc = 1; // 1 == locked
   END_ATOMIC();
   return oldValue;
}
```

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- Atomic operations can be implemented other ways
 - e.g., "test-and-set"

Test-and-set can also be used to implement a spin-lock mutex

```
int testAndSet(int *loc) {
  int oldValue;
  START ATOMIC();
  oldValue = *loc;
  *loc = 1; // 1 == locked
  END ATOMIC();
  return oldValue;
void lock(int *loc) {
  while (testAndSet(loc) == 1);
void unlock(int *loc) {
  START ATOMIC();
  *loc = 0;
  END ATOMIC();
```

 compareAndSwap() provides a more general solution than testAndSet()

```
int testAndSet(int *loc) {
  int oldValue;
  START ATOMIC();
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  END ATOMIC();
  return oldValue;
int compareAndSwap (int *loc,
                    int expected,
                    int updated) {
  START ATOMIC();
  int oldValue = *loc;
  if (oldValue == expected)
     *loc = updated;
  END ATOMIC();
  return oldValue:
```

- compareAndSwap() provides a more general solution than testAndSet()
 - e.g., it can set the value to something other than 1 or 0

```
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  START ATOMIC();
  oldValue = *loc;
  *loc = 1; // 1 == locked
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int compareAndSwap(int *loc,
                    int expected,
                    int updated) {
  START ATOMIC();
  int oldValue = *loc;
  if (oldValue == expected)
     *loc = updated;
  END ATOMIC();
  return oldValue:
```

 One "human" known use of atomic operations is a Star Trek transporter



See en.wikipedia.org/wiki/Transporter_(Star_Trek)

- One "human" known use of atomic operations is a Star Trek transporter
 - Converts a person/object into an energy pattern & "beams" them to a destination where they're converted back into matter



- One "human" known use of atomic operations is a Star Trek transporter
 - Converts a person/object into an energy pattern & "beams" them to a destination where they're converted back into matter
 - This process must occur atomically or a horrible accident will occur!



 Another "human" known use of atomic operations is "apparition" in Harry Potter

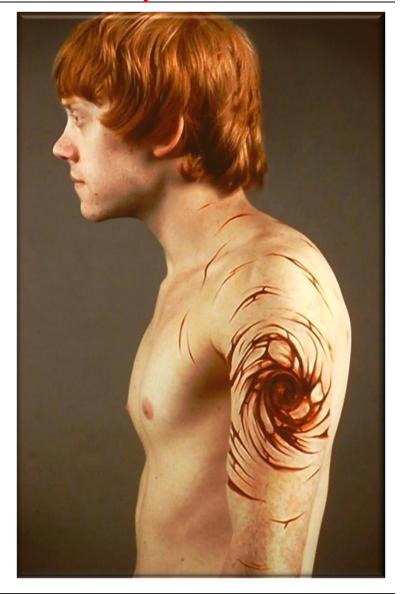


See harrypotter.fandom.com/wiki/Apparition

- Another "human" known use of atomic operations is "apparition" in Harry Potter
 - If the user focuses properly they disappear from their current location & instantly reappear at the desired location



- Another "human" known use of atomic operations is "apparition" in Harry Potter
 - If the user focuses properly they disappear from their current location & instantly reappear at the desired location
 - However, "spinching" occurs if a wizard or witch fails to apparate atomically!



See harrypotter.fandom.com/wiki/Splinching

End of Atomic Classes & Operations: Introduction